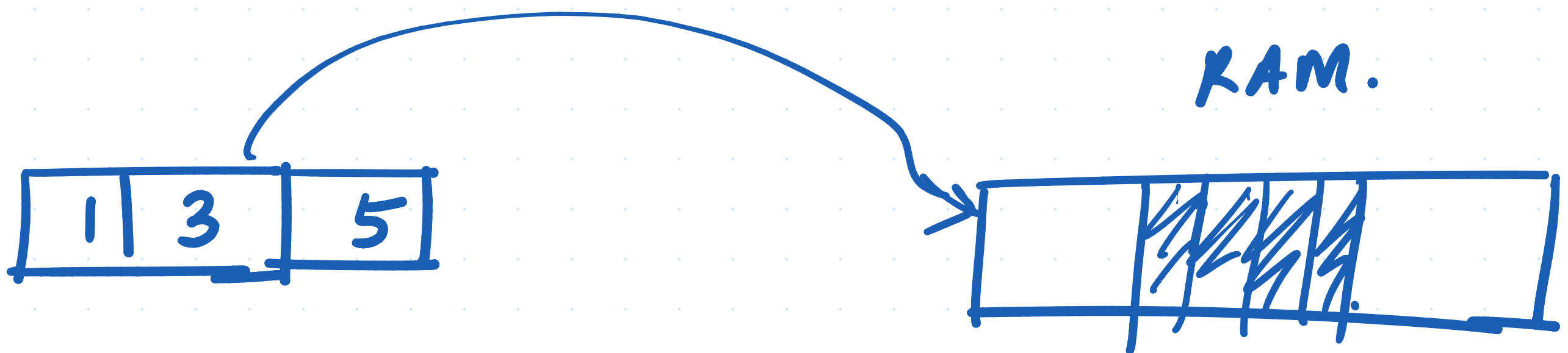


16/08/25.

Arrays / Data Structure.



Measured in Bytes

8 GB of RAM.

Byte $\rightarrow$ 8 bits.

$10^9 \rightarrow$ Bytes

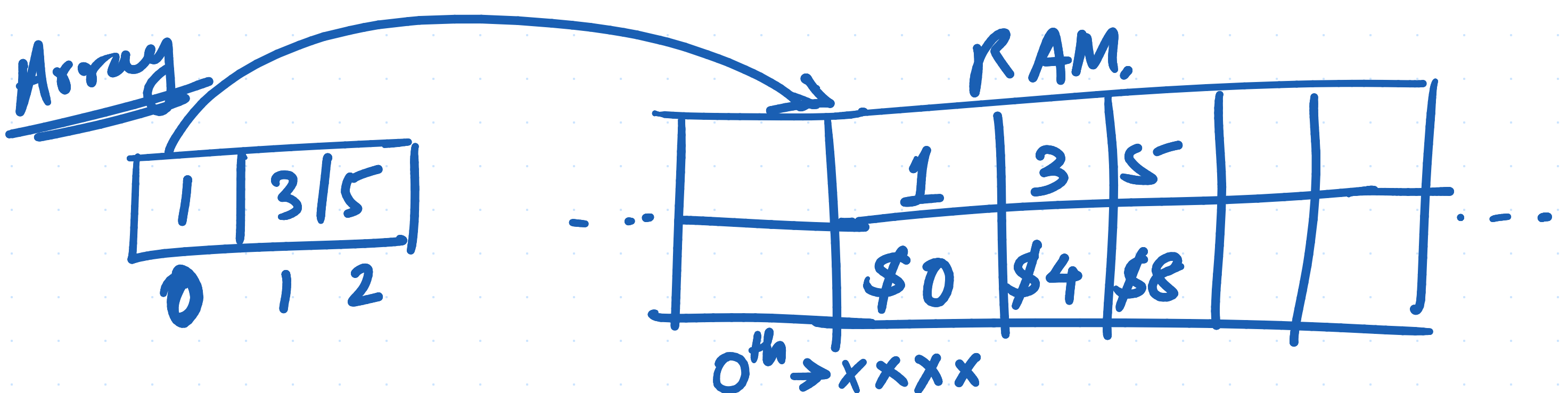
Bit $\rightarrow$ a portion that can store 0/1.

Integer $\rightarrow$ 4 byte.

? Bits? $\rightarrow$ 32 Bits.

1. $\rightarrow$ integer.

0 0 0 - - - - 0 1
 $\leftarrow$ 31 0's $\rightarrow$



$(xxxx + 4) \rightarrow 3 \text{ value.}$

1 byte $\rightarrow$ 1 character

$(xxxx + 8) \rightarrow 5 \text{ value.}$

ASCII value representation.

access.

Read $\rightarrow O(1)$. constant time read

Write $\rightarrow O(1)$ constant time write.

Random Access Memory $\rightarrow$ we can access in constant time, any part of the memory given we have the correct location / address for it.

(Garbage collector).

C, C++
Java.

program.

value	1	3	5		
Address	\$0	\$4	\$8	\$12	\$16

accessed by other / used by others.

C, C++ $\rightarrow$ vectors $\rightarrow$ dynamic.

int arr[10];
 $\downarrow$

Python $\rightarrow$ dynamic array $\rightarrow$ List.

value.	1	5	3	2					
addr	\$0	\$4	\$8	\$12	\$16	\$20	\$24		

$\leftarrow 8 \rightarrow$

inserting in the middle of the array,
we need to shift element, so it will
take $O(n)$ time.

Deletion of first & last $\rightarrow \underline{\underline{O(1)}}$

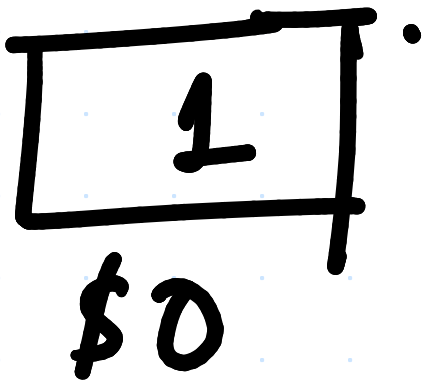
$O(n) \rightarrow$ remove & shift elements of the array

Dynamic Array, how is it maintained in
Python or any kind of Prog Lang.

$\rightarrow$ No size is mentioned.

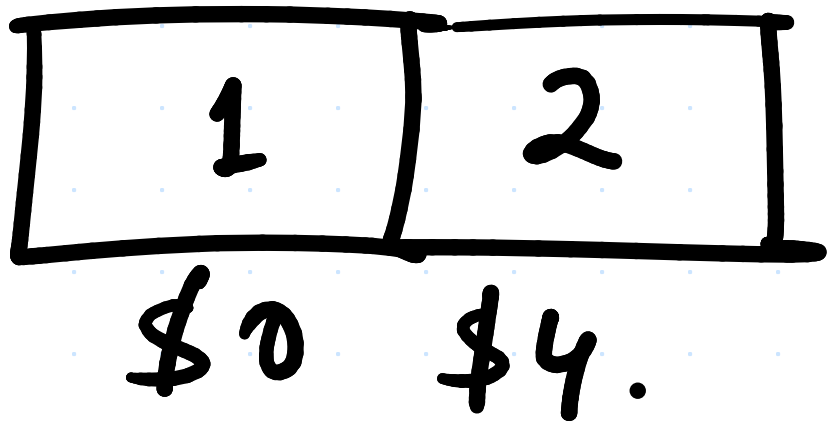
a = [] ; list.

#1



$2 > 1$

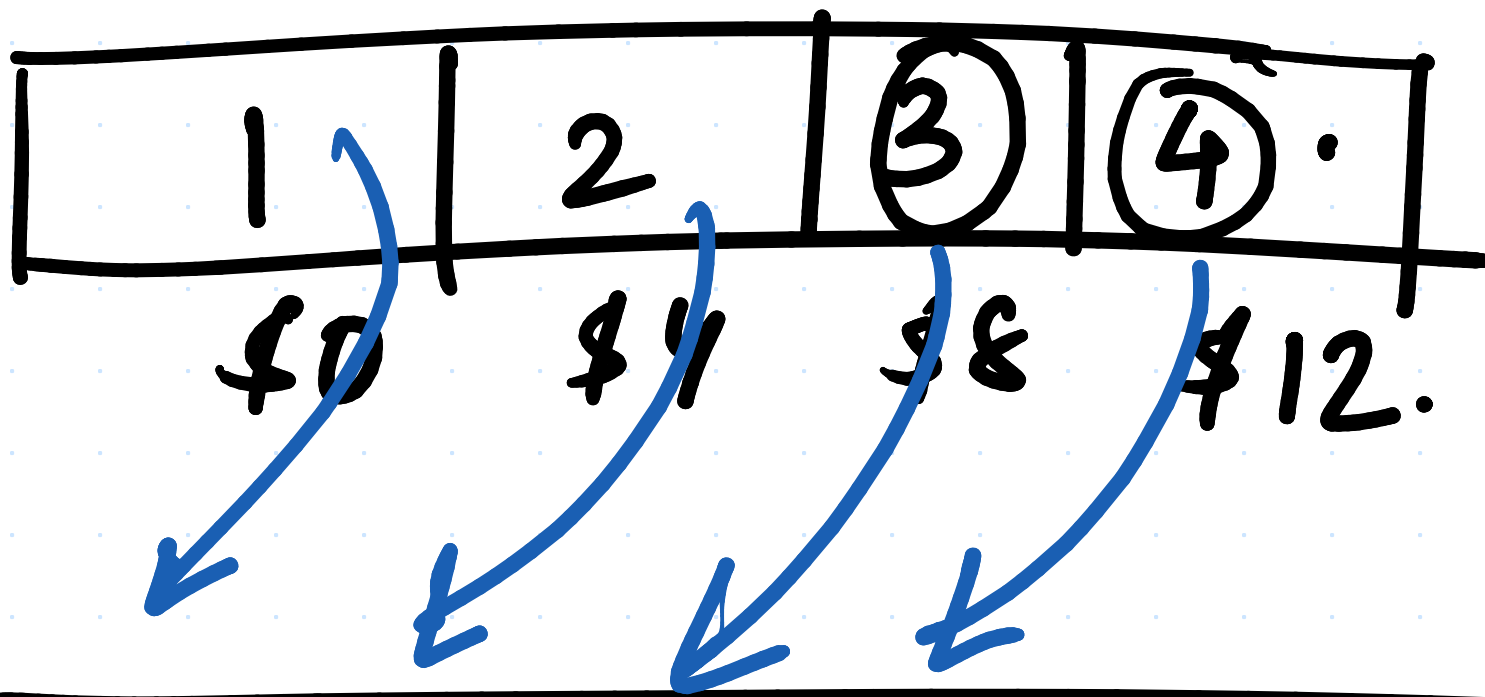
#2



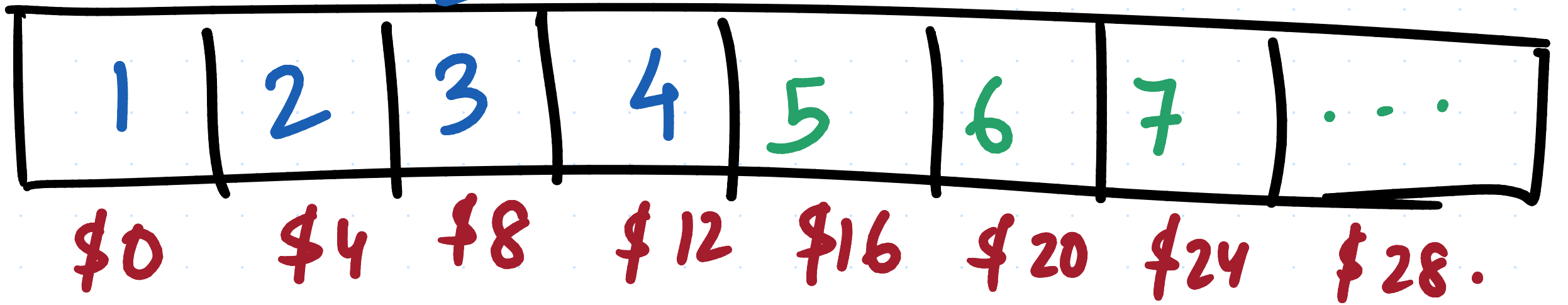
$4 > 2 + 1$

$8 > 4 + 2 + 1$

#4



#8

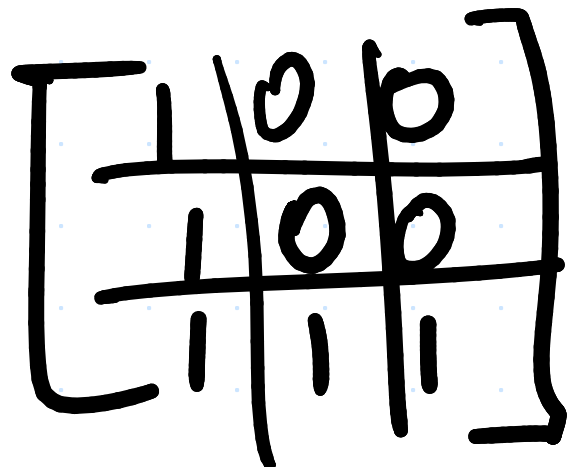
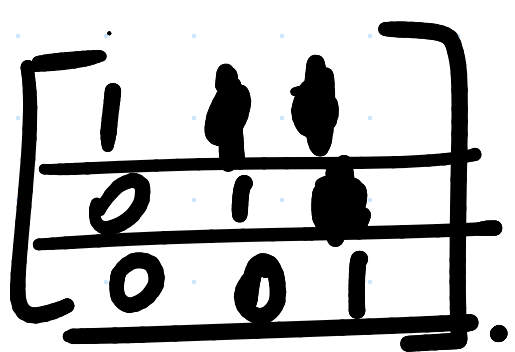
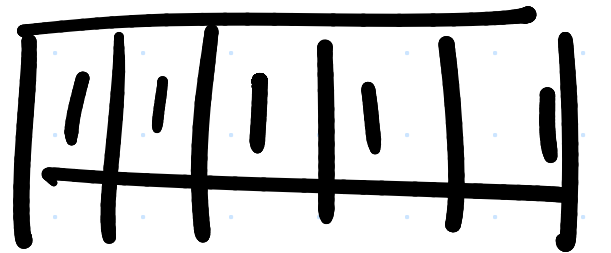


Tradeoff b/w time & space.

Optimal $\rightarrow$ 2x.

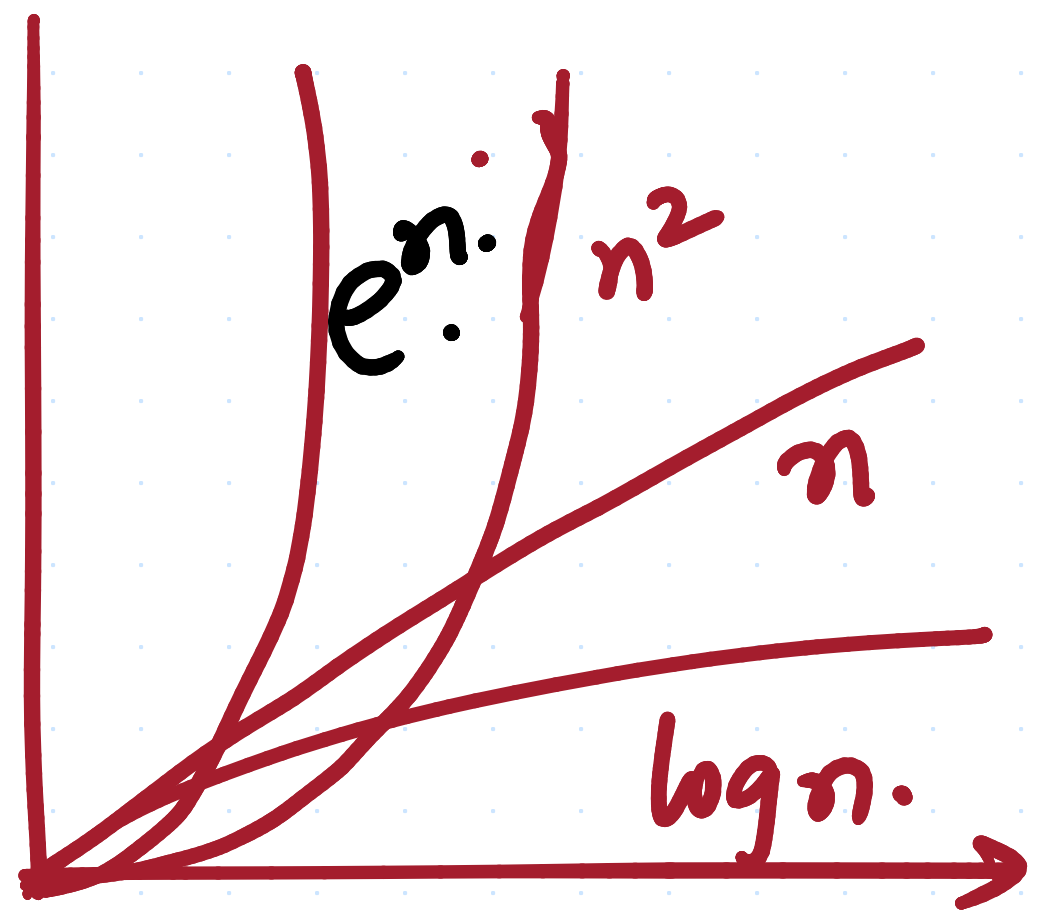
$\rightarrow$ [100]

[200]
[400]



Mostly $O(n)$ & $O(\log n)$

↓
type of Algo.

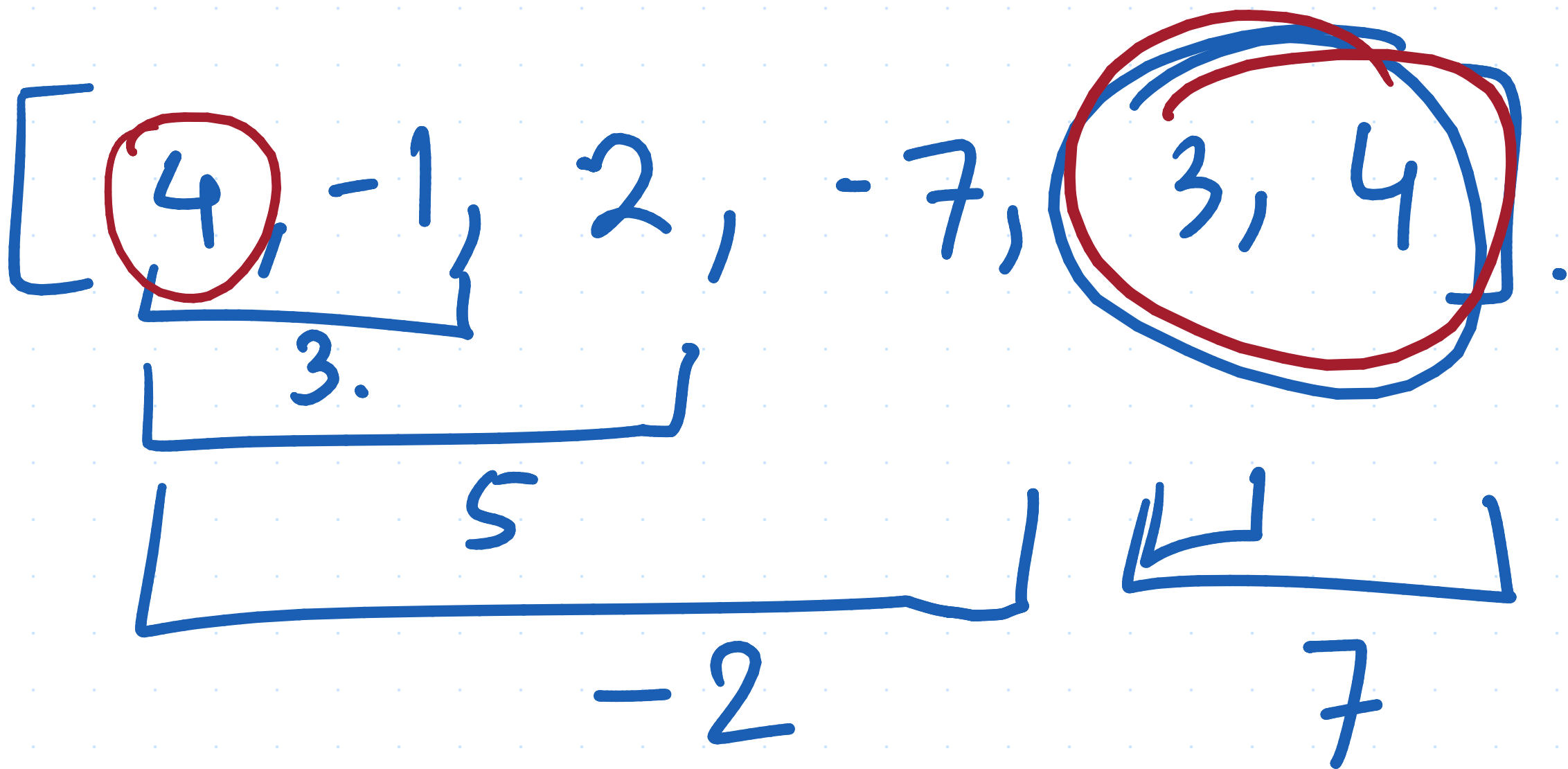


Brute force $\rightarrow O(n^4)$.

Kadane's Algorithm

(2ptr method)

Find a non-empty subarray with the largest sum.



o/p $\rightarrow 7$.

```
def bruteForce (nums)
```

```
    maximum = nums[0]
```

```
    for i in range (len(nums)):
```

```
        cursum = 0.
```

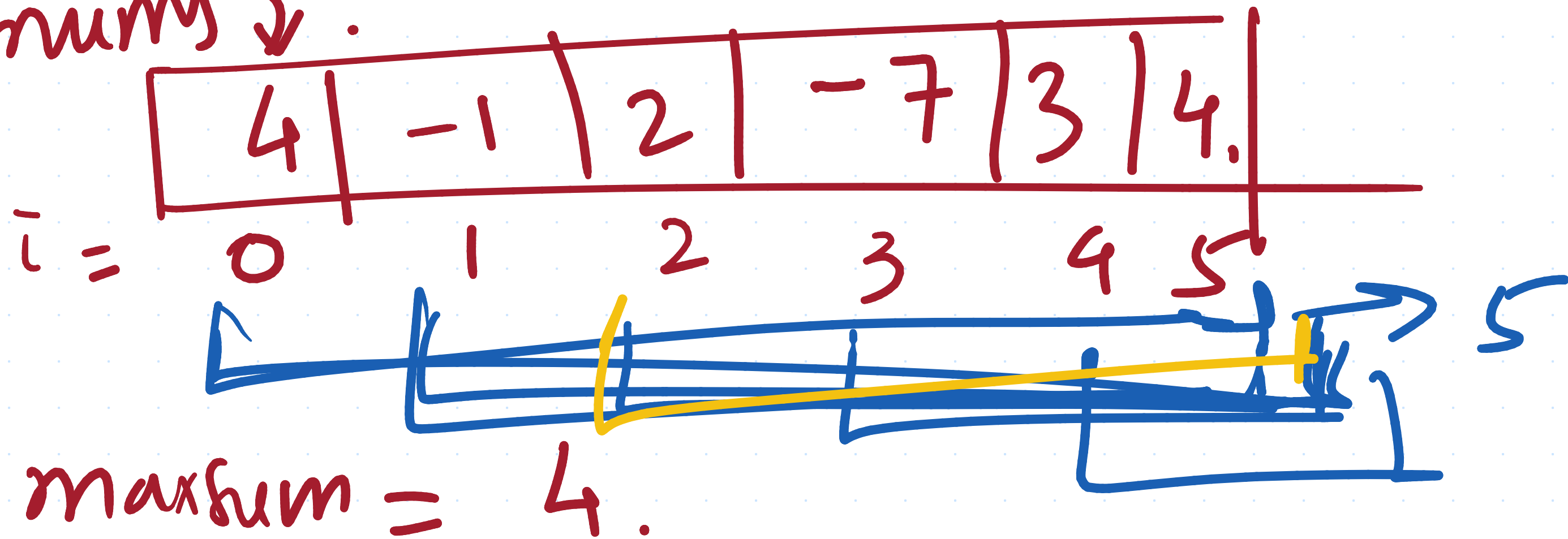
```
        for j in range (i, len(nums)):
```

```
            cursum += nums[j]
```

```
            maximum = max(maximum, cursum).
```

```
    return maximum.
```

nums →



`cursum = 0 4 3 5 -2 1 5.`

`maximum = 4, 4, 4, 5, 5, 5`

current = 0 -1 1 -6 -2 1

maxima = 5 5.5 5 5 5 5



def Kadanes (nums):

maxsum = nums[0]

cursum = 0

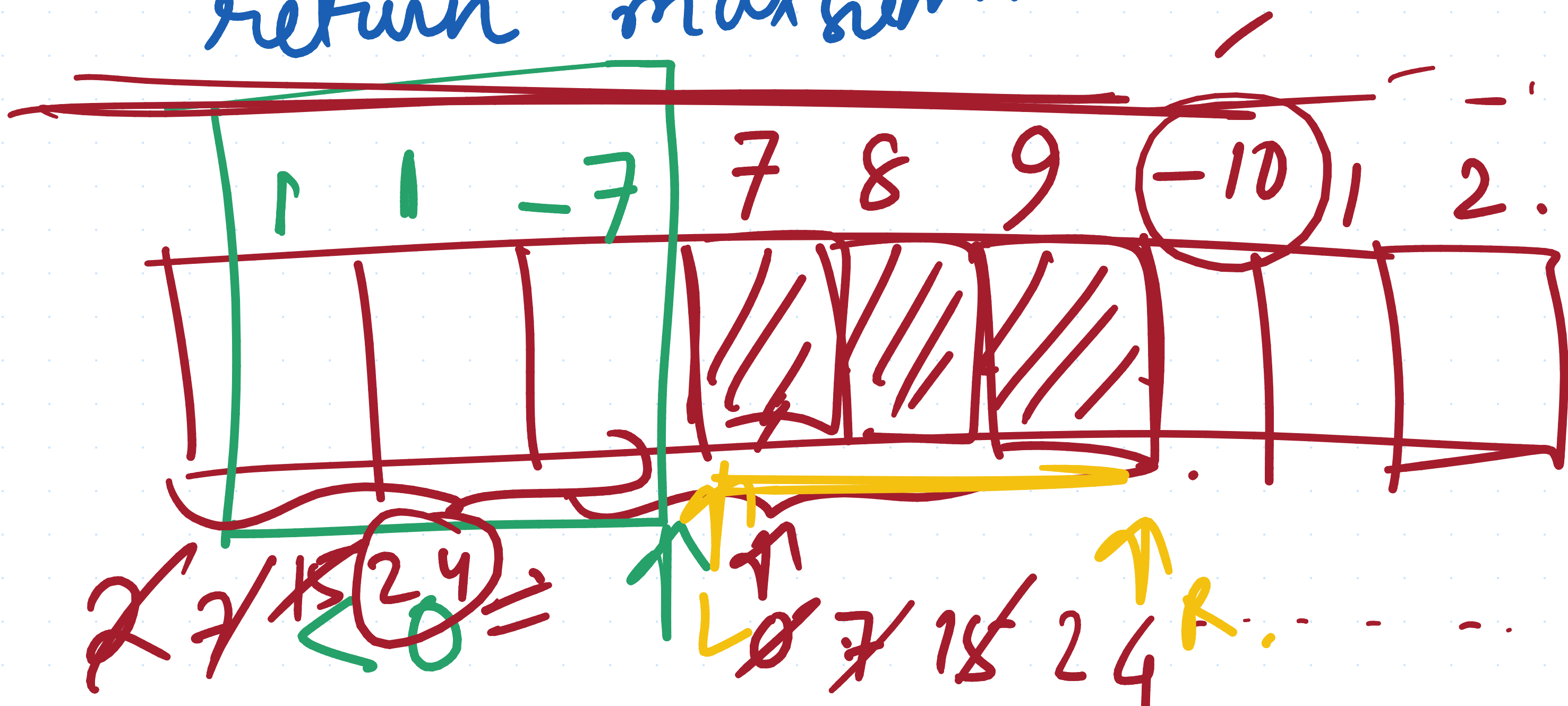
for n in nums: $O(n)$

1 cursum = max(cursum, 0)

1 cursum += n

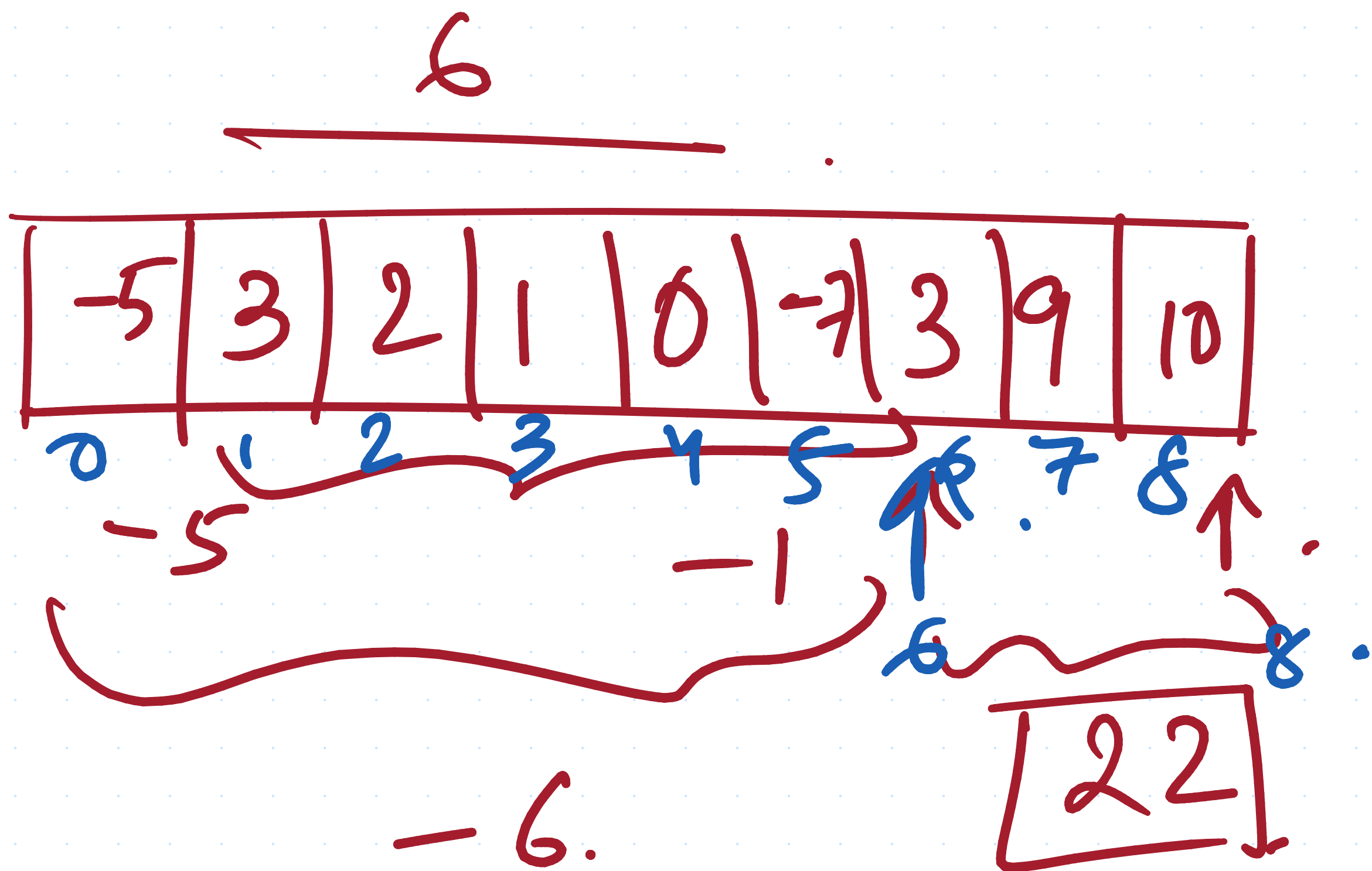
1 maxsum = max(maxsum, cursum)

return maxsum.



sliding window Version of Kadane's Algorithm.

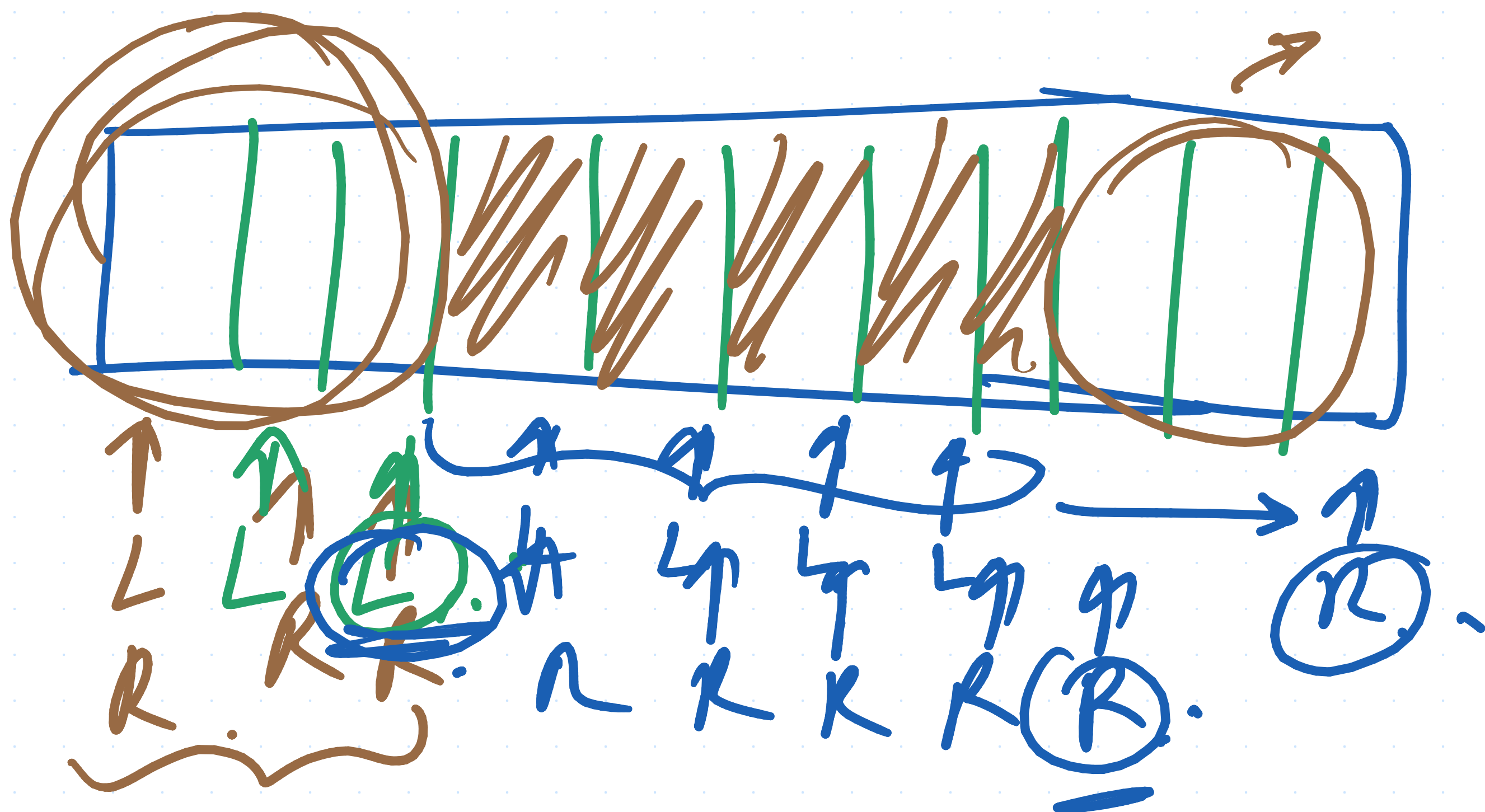
Q: Return the left & right index of the max-subarray sum, assuming there is exactly one result.



```

def slidingWindow(nums):
    maxSum = nums[0]
    curSum = 0
    maxL, maxR = 0, 0
    L = 0
    for R in range(len(nums)):
        if curSum < 0:
            curSum = 0; L = R
        curSum += nums[R]
        if curSum > maxSum:
            maxSum = curSum
            maxL, maxR = L, R
    return [maxL, maxR]

```



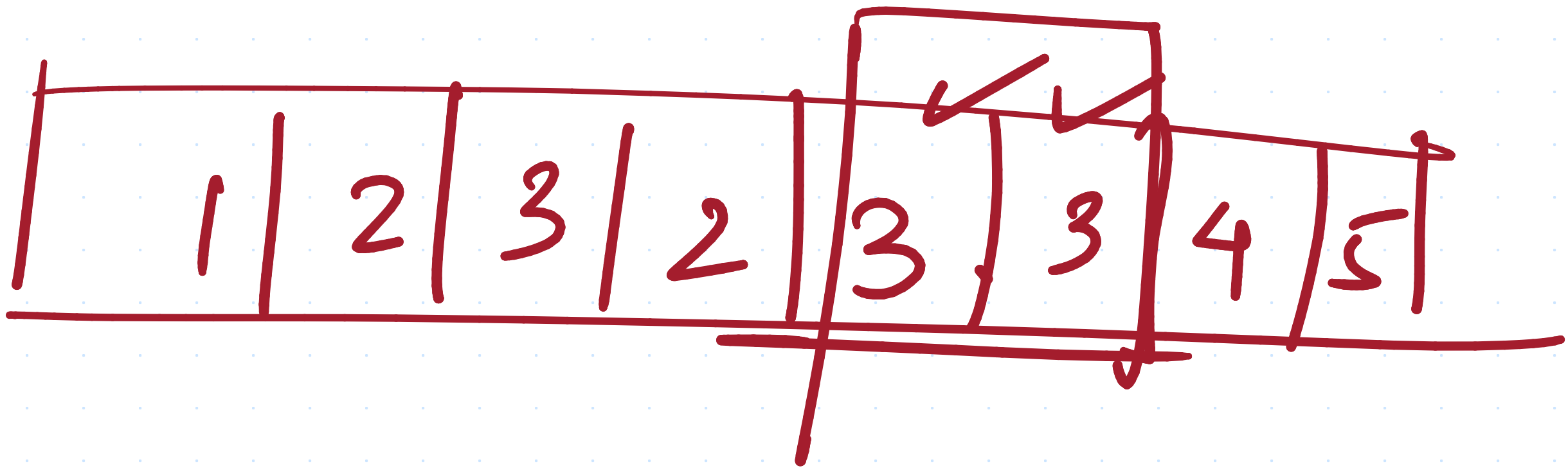
-ve

$$\max L = L;$$

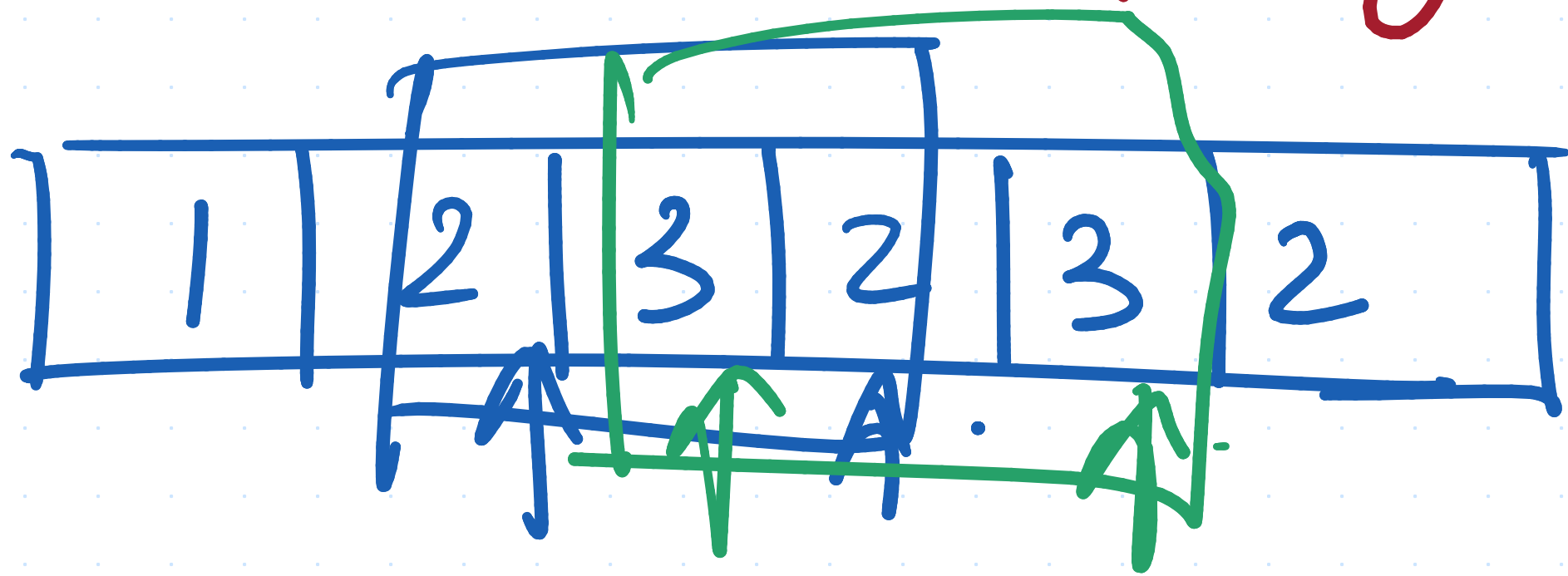
$$\max R = R.$$

Sliding Window

Given an array, return true if there are two elements within a window of size k that are equal.



window of size $k=2$.

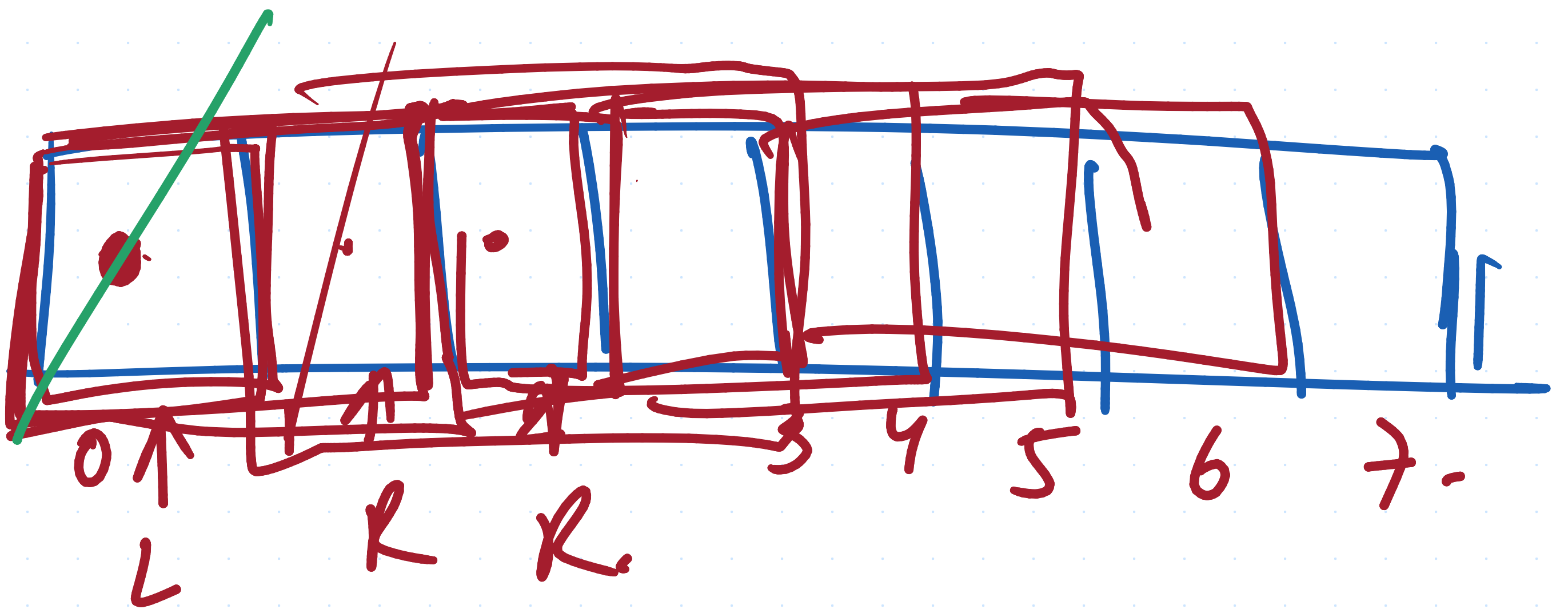
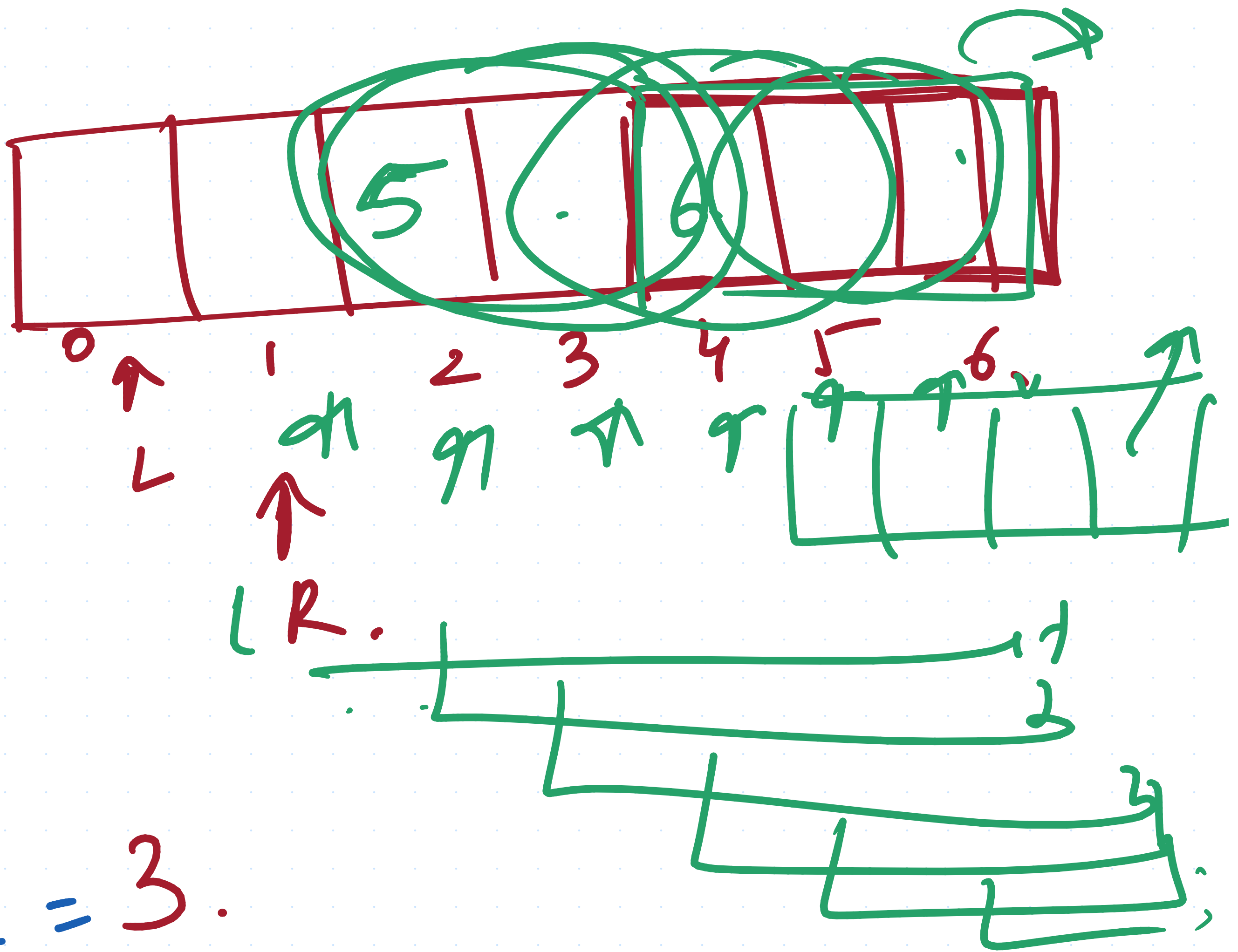


$k=3$.

2

true.

```
def checkDupBF(nums, k):
    for L in range(len(nums)):
        for R in range(L+1,
            min(len(nums), L+k)):
            if nums[L] == nums[R]:
                return True
    return False
```



Sliding window.

$O(1) \rightarrow$ hash set.

$O(n)$ complexity:

```
def closeDuplicates(nums, k):
```

```
    window = set()
```

```
    L = 0
```

```
    for R in range(len(nums)):
```

```
        if  $R - L + 1 > k$ :
```

```
            window.remove(nums[L])
```

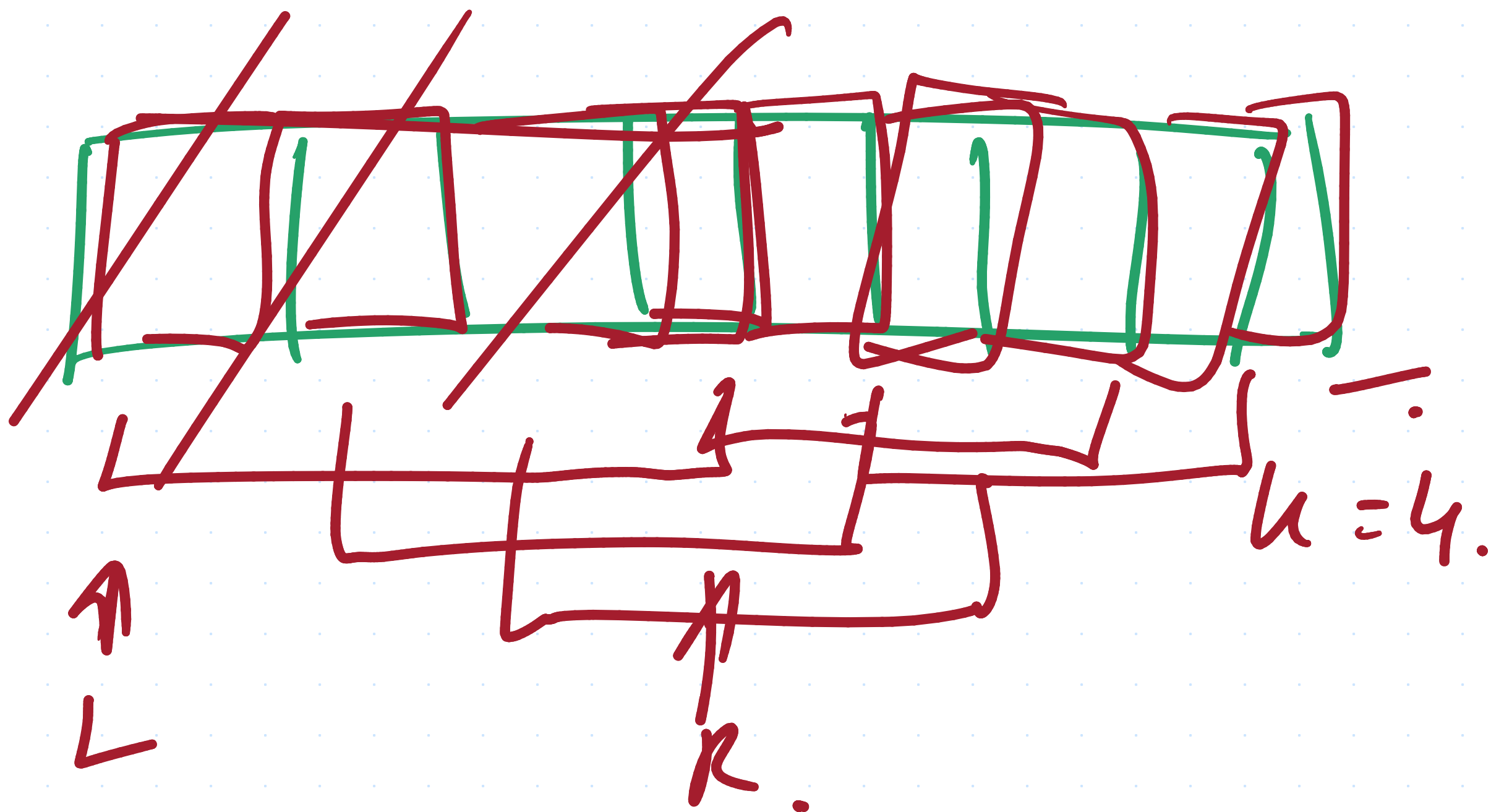
```
             $L += 1$ 
```

```
        if nums[R] in window:
```

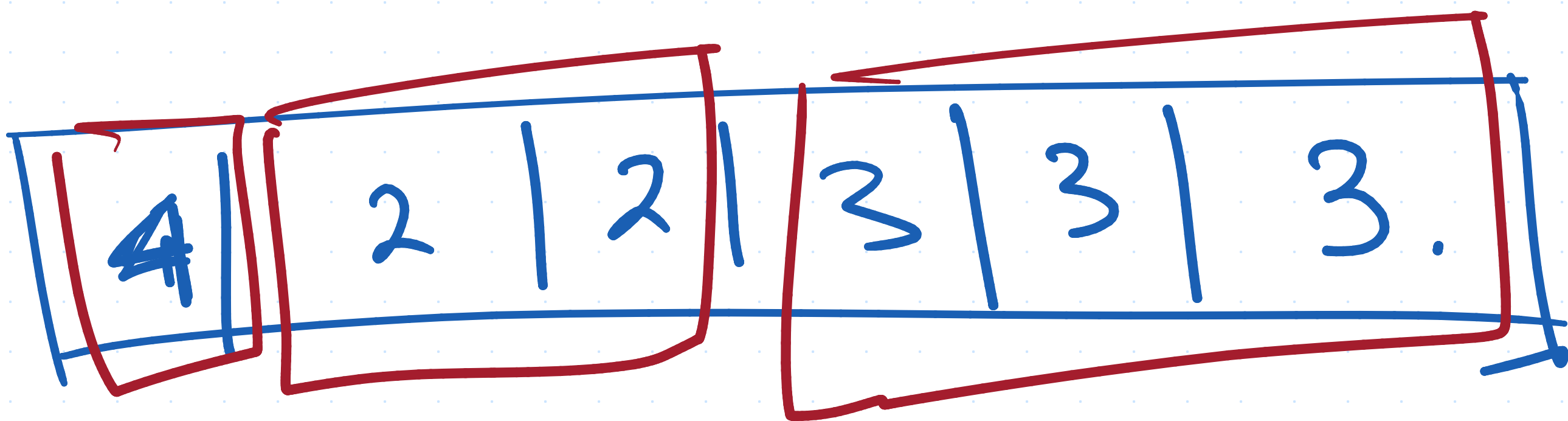
```
            return True
```

```
        window.add(nums[R])
```

```
    return False.
```



Sliding Window with variable length.



• Find the length of the longest subarray with the same value in each position.

$O(n)$

```
def longestSubarray (nums):
```

```
    length = 0; L = 0
```

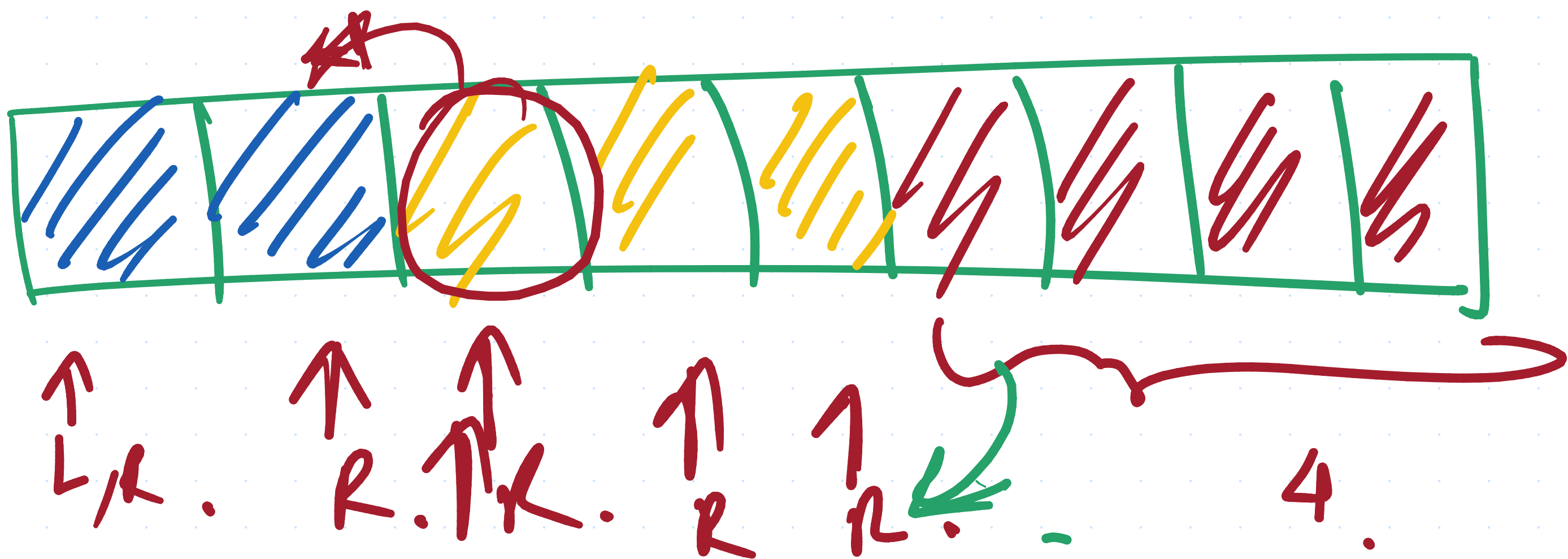
```
    for R in range (len(nums)):
```

```
        if nums[L] != nums[R]:
```

```
            L = R.
```

```
        length = max(length, R-L+1)
```

```
    return length
```



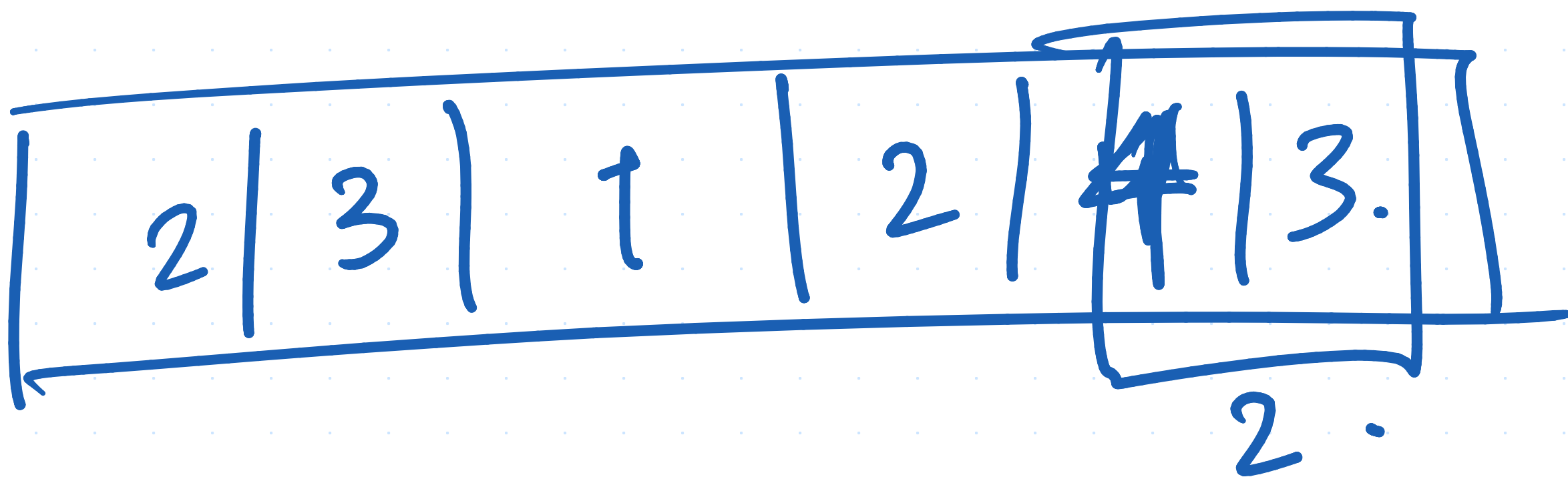
L

L, R, ..., R.

$$= \underline{R - L + 1}$$

find the min length subarray
where the sum is greater
than or equal to the target.

Assume all values are +ve.



Target = 6.

```
def shortestSubarray(nums, target):
```

```
    L, total = 0, 0
```

```
    length = float("inf") #  $\infty$ 
```

```
    for R in range(len(nums)):
```

```
        total += nums[R]
```

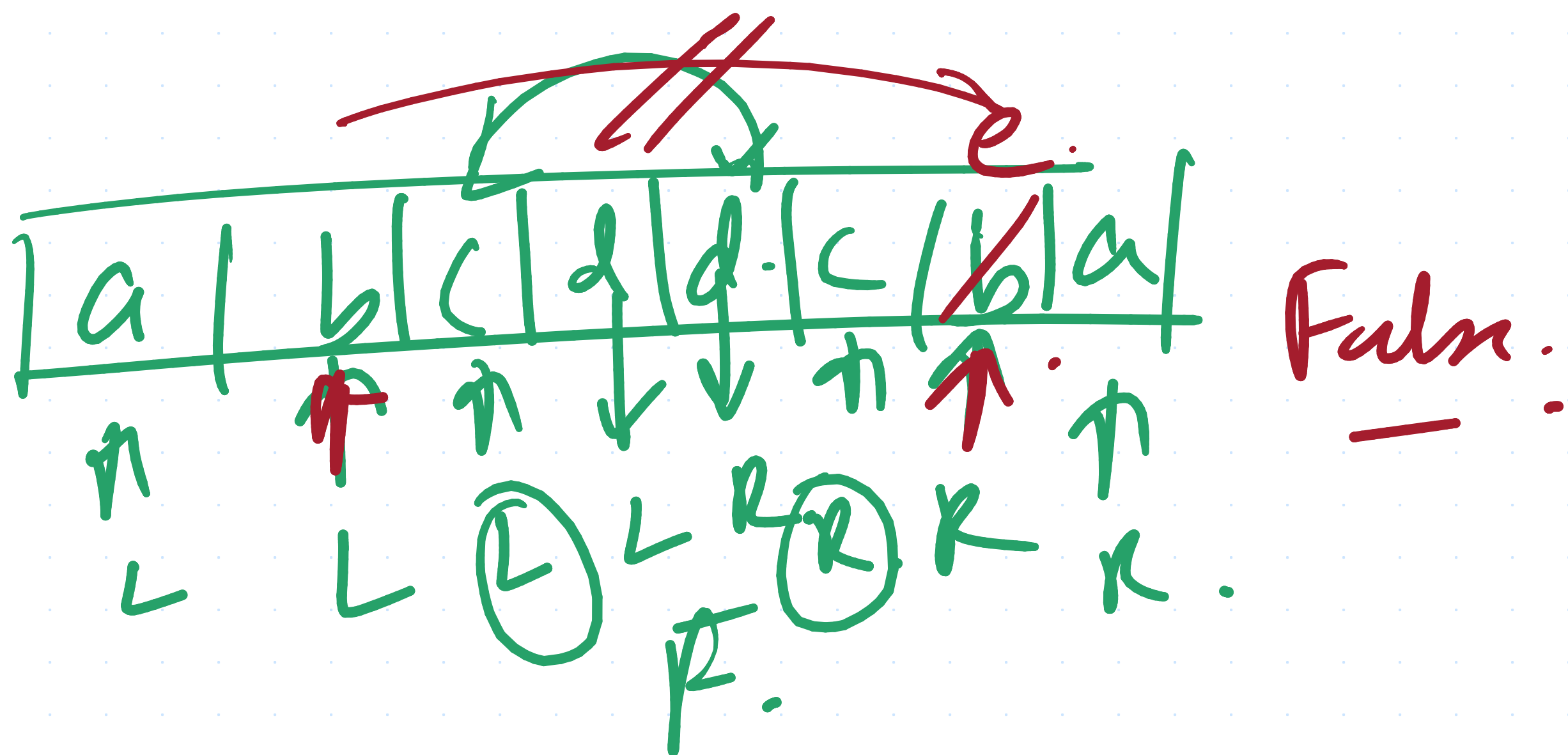
```
        while total >= target:
```

```
            length = min(R - L + 1,  
                          length)
```

```
            total -= nums[L]
```

```
            L += 1
```

```
    return 0 if length == float("inf")
```

check if an array is palindrome.

def isPal(word):

$L \text{ \& } R = 0, \text{ len(word)} - 1$

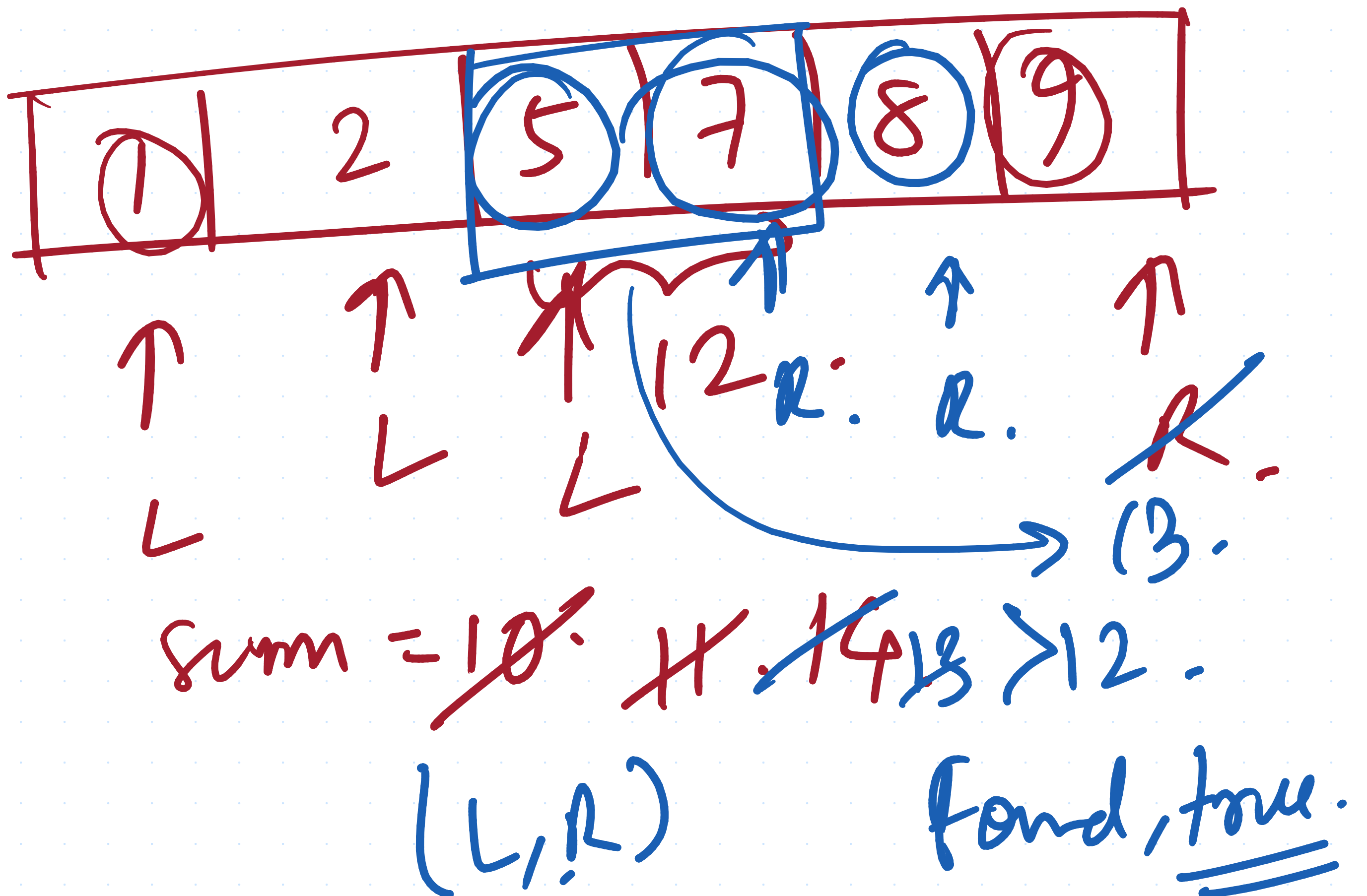
while $L < R$:

if $\text{word}[L] \neq \text{word}[R]$:
return False.

$L += 1; R -= 1$

return True.

Q/ Given a sorted i/p array,
return the two indices of the
two elements which sum up to
the target value. Assume
there is exactly 1 sol.



```
def targetSum (nums, target):
```

```
    L, R = 0, len(nums) - 1
```

```
    while L < R:
```

```
        if nums[L] + nums[R] > target:
```

```
            R -= 1
```

```
        elif nums[L] + nums[R] < target:
```

```
            L += 1
```

```
    else
```

```
        return [L, R]
```
